# Testing Database-Driven Applications: Challenges and Solutions

Gregory M. Kapfhammer Department of Computer Science University of Pittsburgh

Department of Computer Science Allegheny College Mary Lou Soffa Department of Computer Science University of Pittsburgh



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$  esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 1/32

# Outline

- Introduction and Motivation
- Testing Challenges
- Database-Driven Applications
- A Unified Representation
- Test Adequacy Criteria
- Test Suite Execution
- Test Coverage Monitoring
- Conclusions and Resources

# Motivation

The Risks Digest, Volume 22, Issue 64, 2003

#### Jeppesen reports airspace boundary problems

About 350 airspace boundaries contained in Jeppesen NavData are incorrect, the FAA has warned. The error occurred at Jeppesen after a software upgrade when information was pulled from a database containing 20,000 airspace boundaries worldwide for the March NavData update, which takes effect March 20.



# **Testing Challenges**

- Should consider the environment in which software applications execute
- Must test a program and its interaction with a database
- Database-driven application's state space is well-structured, but infinite (Chays et al.)
- Need to show program does not violate database integrity, where *integrity* = *consistency* + *validity* (Motro)
- Must locate program and database coupling points that vary in granularity



# **Testing Challenges**

- The structured query language's (SQL) data manipulation language (DML) and data definition language (DDL) have different interaction characteristics
- Database state changes cause modifications to the program representation
- Different kinds of test suites require different techniques for managing database state during testing



 $\mathcal D$ atabase dr $\mathcal I$ ven  $\mathcal A$ pplication  $\mathcal T$ esting t $\mathcal O$ ol  $\mathcal M$ odule $\mathcal S$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 5/32

# **Testing Challenges**

The many testing challenges include, but are not limited to, the following:

- Unified program representation
- Family of test adequacy criteria
- Efficient test coverage monitoring techinques
- Intelligent approaches to test suite execution



# **Database-Driven Applications**







 Program P interacts with two relational databases

 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 7/32



 A program can interact with a database at different levels of granularity



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 8/32



I A T O M S

 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$  esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 8/32





 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 8/32









### **Database Interaction Points**

- → Database interaction point  $I_r \in I$  corresponds to the execution of a SQL DML statement
- Consider a simplified version of SQL and ignore SQL
   DDL statements (for the moment)
- Interaction points are normally encoded within Java programs as dynamically constructed Strings
- → select uses D<sub>k</sub>, delete defines D<sub>k</sub>, insert defines D<sub>k</sub>,
   update defines and/or uses D<sub>k</sub>



### **Database Interaction Points (DML)**

```
select A_1, A_2, \ldots, A_q
from r_1, r_2, \ldots, r_m
where Q
```

delete from rwhere Q

insert into  $r(A_1, A_2, ..., A_q)$  update rvalues $(v_1, v_2, ..., v_q)$  set  $A_l = F(A_l)$ where Q



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 10/32

# **Refined Database-Driven Application**



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 11/32

### **Test Adequacy Criteria**

- *P* violates a database  $D_k$ 's validity when it:
  - → (1-v) inserts entities into D<sub>k</sub> that do not reflect real world
- *P* violates a database  $D_k$ 's completeness when it:
  - → (1-c) deletes entities from D<sub>k</sub> that still reflect real world
- → In order to verify (1-v) and (1-c), T must cause P to define and then use entities within  $D_1, \ldots, D_n$ !



### **Data Flow Information**

- Interaction point: ``UPDATE UserInfo SET
   acct\_lock=1 WHERE card\_number='' +
   card\_number + ``;'';
  - Database Level: *define(BankDB)*
  - Attribute Level: define(acct\_lock) and use(card\_number)
- Data fbw information varies with respect to the granularity of the database interaction



### **Database Entities**

**UserInfo** 

card_number	pin_number	user_name	acct_lock
1	32142	Brian Zorman	0
2	41601	Robert Roos	0
3	45322	Marcus Bittman	0
4	56471	Geoffrey Arnold	0

$$A_{v}(I_{r}) = \{ 1, 32142, \ldots, \text{Geoffrey Arnold}, 0 \}$$



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 14/32

# **The DICFG: A Unified Representation**



- "Database-enhanced"
   CFG for lockAccount
- Define temporaries to represent the program's interaction at the levels of
   database and attribute



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 15/32

# **The DICFG: A Unified Representation**



- Database interaction
   graphs (DIGs) are
   placed before interaction
   point I<sub>r</sub>
- Multiple DIGs can be integrated into a single CFG
- String at I<sub>r</sub> is
   determined in a
   control-fbw sensitive
   fashion

# **Test Adequacy Criteria**



- Database interaction
   association (DIA) involves the def and use of a database
   entity
- DIAs can be located in the DICFG with data flow analysis
  - all-database-DUs requires
     tests to exercise all DIAs for all
     of the accessed databases



# **Generating Test Requirements**



Measured time and space overhead when computing family of test adequacy criteria

- Modifi ed ATM and mp3cd to contain appropriate database interaction points
- Soot 1.2.5 to calculate intraprocedural associations
- GNU/Linux workstation with kernel
   2.4.18-smp and dual 1 GHz Pentium III
   Xeon processors



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 17/32

### **Counting Associations and Definitions**



 DIAs at attribute value level represent 16.8% of mp3cd's and 9.6% of ATM's total number of intraprocedural associations

### **Measuring Time Overhead**



 Computing DIAs at the attribute value level incurs no more than a 5 second time overhead

#### **Measuring Average Space Overhead**



mp3cd shows a more marked increase in the average number of nodes and edges than ATM

### **Measuring Maximum Space Overhead**



mp3cd shows a signifi cantly greater maximum space overhead than ATM

# **Automatic Representation Construction**

- Manual construction of DICFGs is not practical
- Use extension of BRICS Java String Analyzer (JSA) to determine content of String at I<sub>r</sub>
- Per-class analysis is inter-procedural and control flow sensitive
- Conservative analysis might determine that all database entities are accessed
- Include coverage monitoring instrumentation to track
   DIGs that are covered during test suite execution



# **Tracking Covered DIGs and DIAs**





 DIA coverage can be tracked by recording which DIGs within a DICFG were executed during testing

 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 23/32

# **Types of Test Suites**









**Non-restricted** 

**Partially Independent** 

### **Test Suite Execution**

- → Independent test suites can be executed by using provided setup code to ensure that all  $\Delta_{\gamma} = \Delta_0$
- Non-restricted test suites simply allow state to accrue
- Partially independent test suites must return to Δ<sub>ε</sub> after
   T<sub>ε</sub> is executed by :
  - 1. Re-executing all SQL statements that resulted in the creation of  $\Delta_{\varepsilon}$
  - 2. Creating a compensating transaction to undo the SQL statements executed by each test after  $T_{\varepsilon}$



### **Representation Extension**

- The execution of a SQL insert during testing requires the re-creation of DICFG(s)
- The SQL delete does not require re-creation because we must still determine if deleted entity is ever used
- DICFG re-creation only needed when database interactions are viewed at the record or attribute-value level
- Representation extension ripples to other methods
- DICFGs can be re-constructed after test suite has executed, thus incurring smaller time overhead

 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 26/32

# **Test Coverage Monitoring**

- For each tested method m<sub>i</sub> that interacts with a database and each interaction point I<sub>r</sub> that involves an insert we must:
  - 1. Update the DICFG
  - 2. Re-compute the test requirements
- We can compute the set of covered DIAs by consulting the DIG coverage table
- Test adequacy is : # covered DIAs / # total DIAs



# **Calculating Adequacy**



Test Requirements  $\mathcal{M}_i$ 

DIA	COV?	
<def(e1), use(e1)=""></def(e1),>	$\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{\mathbf{$	
<def(e2), use(e2)=""></def(e2),>		
<def(e3), use(e3)=""></def(e3),>		
<def(e4), use(e4)=""></def(e4),>	>	







$$cov(m_i) = \frac{2}{4} \quad cov(m_j) = \frac{4}{6} \quad cov(T_f) = \frac{6}{10}$$

 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 28/32

### **Related Work**

- Jin and Offutt and Whittaker and Voas have suggested that the environment of a software system is important
- Chan and Cheung transform SQL statements into C code segments
- Chays et al. and Chays and Deng have created the category-partition inspired AGENDA tool suite
- Neufeld et al. and Zhang et al. have proposed techniques for database state generation



Dauo et al. focused on the regression testing of database-driven applications

# Conclusions

- Must test the program's interaction with the database
- Many challenges associated with (1) unified program representation, (2) test adequacy criteria, (3) test coverage monitoring, (4) test suite execution
- The DICFG shows database interactions at varying levels of granularity
- Unique family of test adequacy criteria to detect type (1)
   violations of database validity and completeness



Intraprocedural database interactions can be computed from a DICFG with minimal time and space overhead

# Conclusions

- Test coverage monitoring instrumentation supports the tracking of DIAs executed during testing
- Three types of test suites require different techniques to manage the state of the database
- SQL insert statement causes the re-creation of the representation and re-computation of test requirements
- Data fbw-based test adequacy criteria can serve as the foundation for automatically generating test cases and supporting regression testing



#### Resources

Gregory M. Kapfhammer and Mary Lou Soffa. A Family of Test Adequacy Criteria for Database-Driven Applications. In FSE 2003.

Gregory M. Kapfhammer. Software Testing. CRC Press Computer Science Handbook. June, 2004.

http://cs.allegheny.edu/~gkapfham/research/diatoms/



 $\mathcal{D}$ atabase dr $\mathcal{I}$ ven  $\mathcal{A}$ pplication  $\mathcal{T}$ esting t $\mathcal{O}$ ol  $\mathcal{M}$ odule $\mathcal{S}$ , IBM T.J. Watson Research Center, May 14, 2004 – p. 32/32